

ABI与内联汇编

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本讲概述

i386, IA-32, x86-64, x64,

- 本讲内容
 - 一些背景
 - ABI
 - x86-64 体系结构与汇编语言
 - inline assembly

机器字长的发展

字长 (Word size)

In computing, a word is the natural unit of data used by a particular processor design. The *number of bits in a word* (the word size, word width, or word length)...

- 能直接进行整数/位运算的大小
- 指针的大小 (索引内存的范围)

8 位机 (6502)

- 16 bit PC 寄存器 (64 KiB 寻址能力, KiB 级内存, 无乘法/浮点数)
 - Apple II; Atari 2600; NES; ...



16位机 (8086/8088)

- 我们需要更大的内存！更大的数据宽度！
 - 20 bit 地址线 (两个寄存器)
 - Intel 8086：经典中的经典→x86架构



32 位机 (Intel x86)

- 80386 处理器 4,096 倍的地址空间
 - 4 GiB 内存在 1980s 看起来是非常不可思议的



64 位机 (x86-64; AArch64; RV64; ...)

- 64 位地址空间能索引 17,179,869,184 GiB 内存
 - 我们的服务器有 128 GiB 内存
 - 目前看起来是非常够用的 (PML4)
 - 现在的处理器一般实现 48 bit 物理地址 (256 TiB)

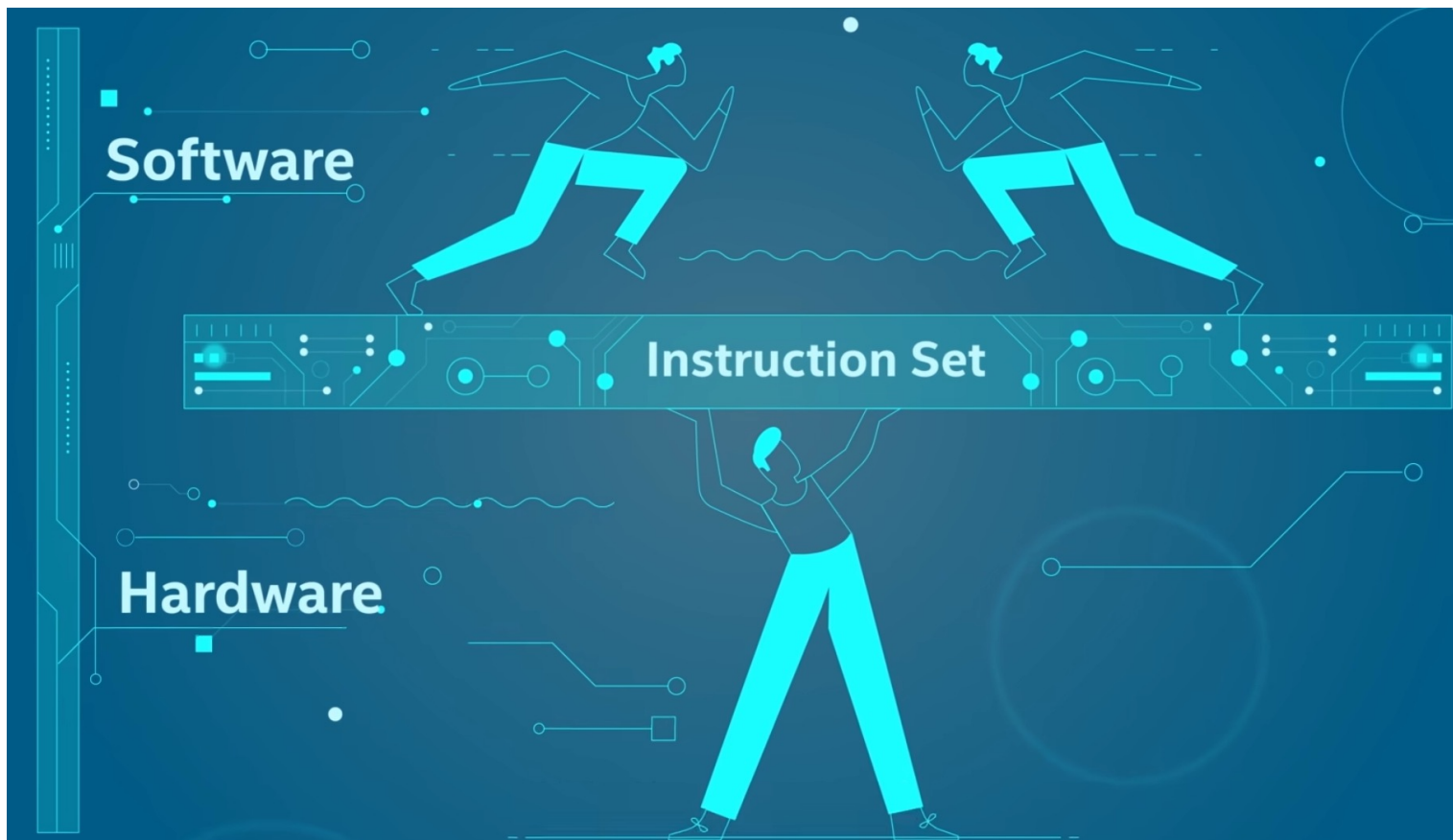


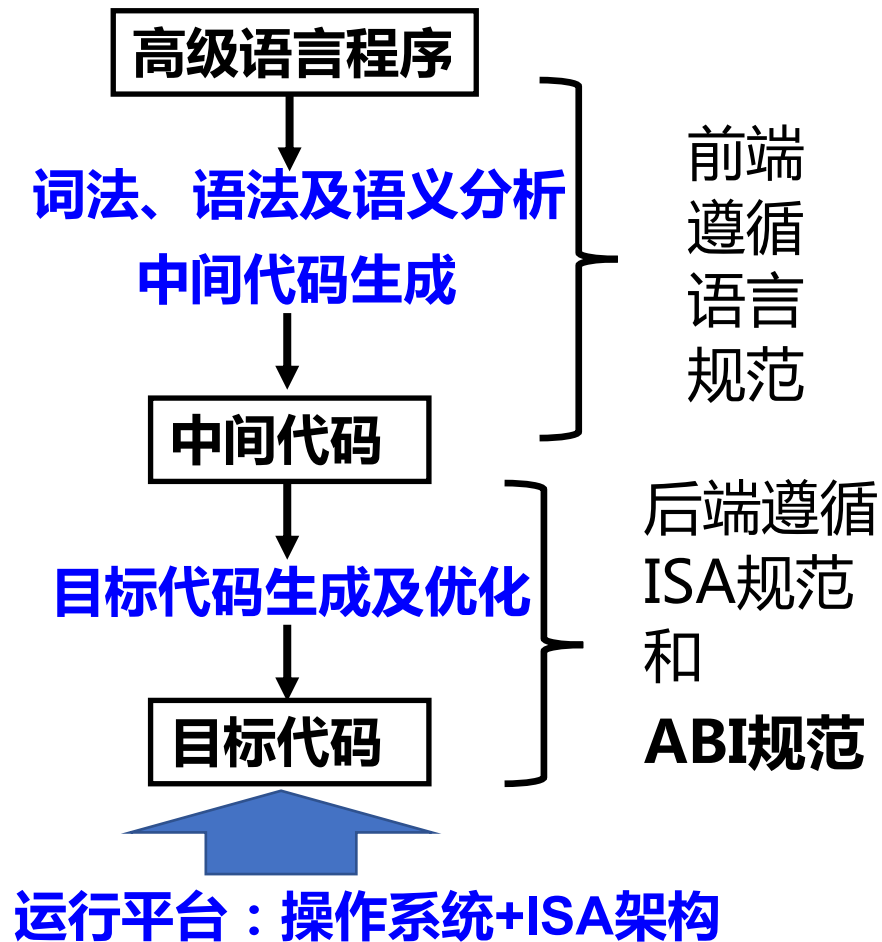
Fun Facts

- `int`类型的长度
 - 8 bit computer: 8 bit
 - 16 bit computer: 16 bit
 - 32 bit computer: 32 bit
 - 64 bit computer: 32 bit
 - JVM (32/64 bit): 32 bit
- 在逻辑世界里描述日常世界 , 2,147,483,647 已经足够大了

处理器架构和指令集的演进

- x86 , x64 , amd64 , ...
- x86 , ARM , MIPS , RISC , ...

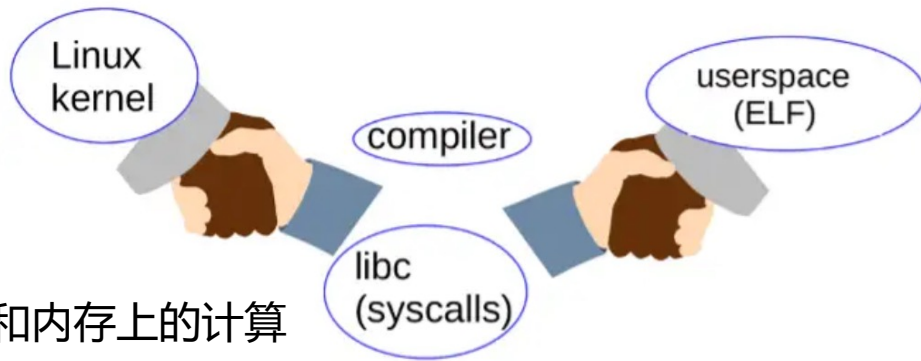




概念复习：ABI

程序的机器级表示

- 程序经历： $.c \rightarrow .o$ (编译)和 $.o \rightarrow a.out$ (链接)
 - 不同版本、不同编译器、不同语言的二进制文件都可以链接
 - 他们需要一个“共同语言”
- ABI (Application Binary Interface)
 - 区别于API (Application Programming Interface)
 - 程序源代码中的规范
 - 约定binary的行为 ([Linux ABI](#))
 - 二进制文件的格式
 - 函数调用、系统调用.....
 - C语言规范只定义了运行时内存和内存上的计算
 - printf都无法实现，必须借助外部的库函数
 - 链接、加载的规范



System V ABI

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Application Binary Interface (ABI)

- 例如：我们熟悉的x86 calling convention
 - cdecl (Linux)
 - stdcall (Win32)
 - 只要遵循标准的函数就可以互相调用

```
C stdio.h  X
D: > gcc > mingw64 > x86_64-w64-mingw32 > include > C stdio.h > fprintf(FILE *__restrict__, const char *__restrict__, ...)

424 #pragma GCC diagnostic push
425 #pragma GCC diagnostic ignored "-Wshadow"
426 #endif
427 __attribute__((__format__ (__MINGW_PRINTF_FORMAT, 2, 3))) __MINGW_ATTRIB_NONNULL(2)
428 int __cdecl fprintf(FILE *__restrict__ _File, const char *__restrict__ _Format, ...);
429 __attribute__((__format__ (__MINGW_PRINTF_FORMAT, 1, 2))) __MINGW_ATTRIB_NONNULL(1)
430 int __cdecl printf(const char *__restrict__ _Format, ...);
431 __attribute__((__format__ (__MINGW_PRINTF_FORMAT, 2, 3))) __MINGW_ATTRIB_NONNULL(2)
432 int __cdecl sprintf(char *__restrict__ _Dest, const char *__restrict__ _Format, ...) __M
433
```

例子：cdecl函数调用

- caller stack frame
 - 所有参数以数组的形式保存在堆栈上（所以才有“反序压栈”）
 - 然后是返回地址
 - 跳转到callee
- callee
 - EAX作为返回值
 - 其他寄存器都是callee save

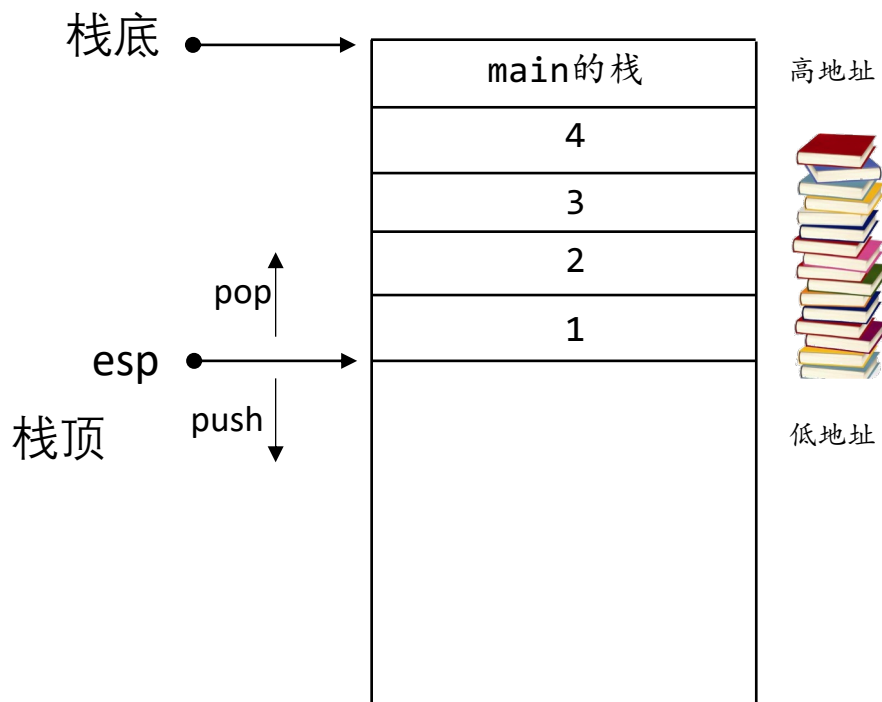
```
void bar(int *);  
int foo(int x) {  
    bar(&x);  
    return x;  
}
```

cdecl函数调用惯例

```
int _cdecl foo(int m1, int m2, int m3, int m4);
```

```
int foo(int m1, int m2, int m3, int m4) {  
    return m1 + m2 + m3 + m4;  
}
```

```
int main() { foo(1,2,3,4);} 
```

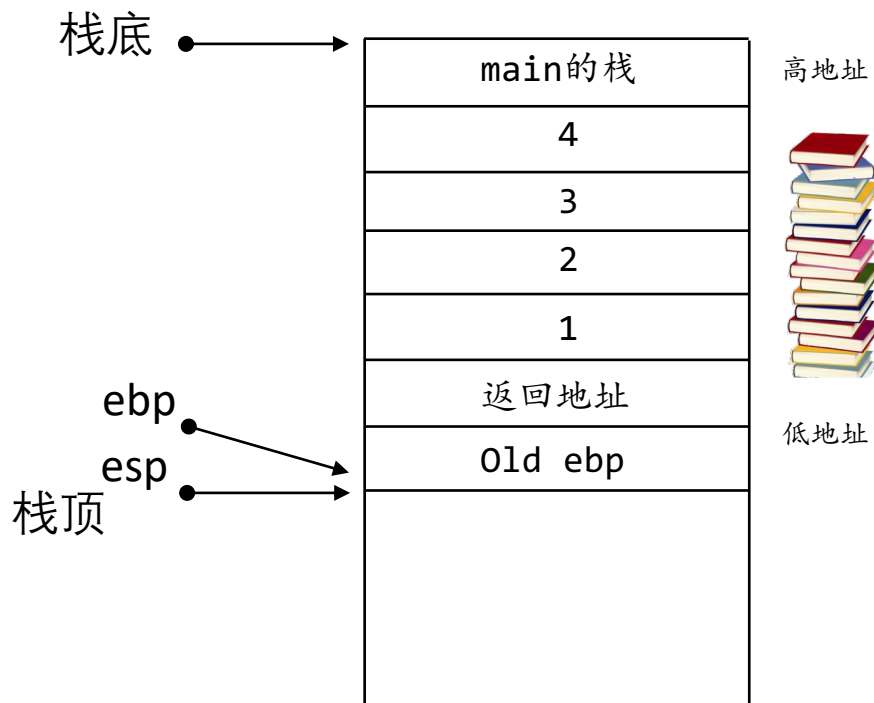


cdecl函数调用惯例

```
int _cdecl foo(int m1, int m2, int m3, int m4);
```

```
int foo(int m1, int m2, int m3, int m4) {  
    return m1 + m2 + m3 + m4;  
}
```

```
int main() { foo(1,2,3,4);}
```



Disassembly of section `.text`:

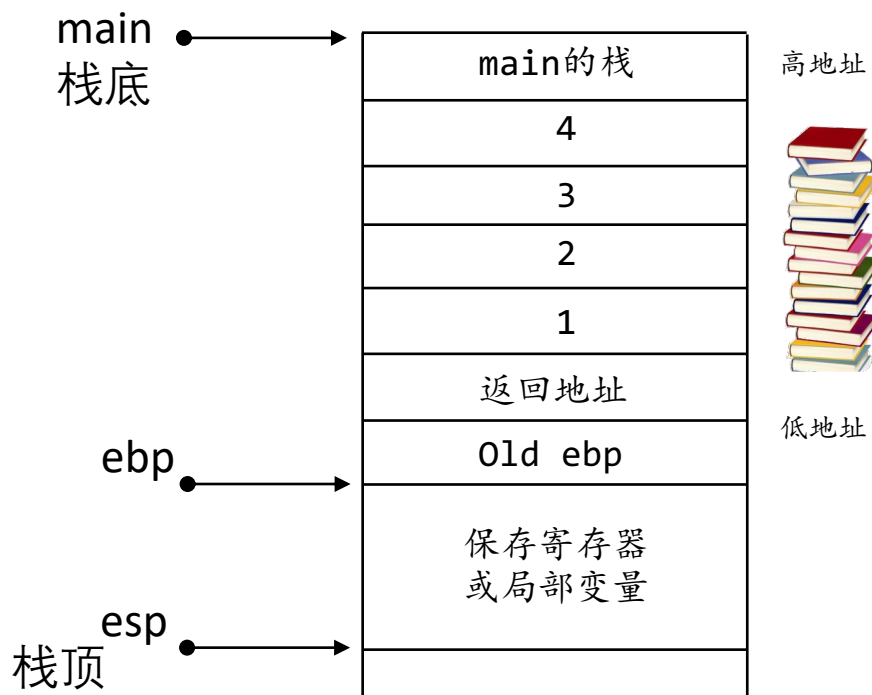
```
00000000 <foo>:  
  0:  f3 0f 1e fb      endbr32  
  4:  55               push    %ebp  
  5:  89 e5            mov     %esp,%ebp  
  7:  e8 fc ff ff ff   call   8 <foo+0x8>  
  c:  05 01 00 00 00   add     $0x1,%eax  
 11:  8b 55 08         mov     0x8(%ebp),%edx  
 14:  8b 45 0c         mov     0xc(%ebp),%eax  
 17:  01 c2           add     %eax,%edx  
 19:  8b 45 10         mov     0x10(%ebp),%eax  
 1c:  01 c2           add     %eax,%edx  
 1e:  8b 45 14         mov     0x14(%ebp),%eax  
 21:  01 d0           add     %edx,%eax  
 23:  5d              pop     %ebp  
 24:  c3              ret  
  
00000025 <main>:  
 25:  f3 0f 1e fb      endbr32  
 29:  55               push    %ebp  
 2a:  89 e5            mov     %esp,%ebp  
 2c:  e8 fc ff ff ff   call   2d <main+0x8>  
 31:  05 01 00 00 00   add     $0x1,%eax  
 36:  6a 04           push    $0x4  
 38:  6a 03           push    $0x3  
 3a:  6a 02           push    $0x2  
 3c:  6a 01           push    $0x1  
 3e:  e8 fc ff ff ff   call   3f <main+0x1a>  
 43:  83 c4 10         add     $0x10,%esp  
 46:  b8 00 00 00 00   mov     $0x0,%eax  
 4b:  c9              leave   %eax  
 4c:  c3              ret
```

cdecl函数调用惯例

```
int _cdecl foo(int m1, int m2, int m3, int m4);
```

```
int foo(int m1, int m2, int m3, int m4) {  
    return m1 + m2 + m3 + m4;  
}
```

```
int main() { foo(1,2,3,4);}
```



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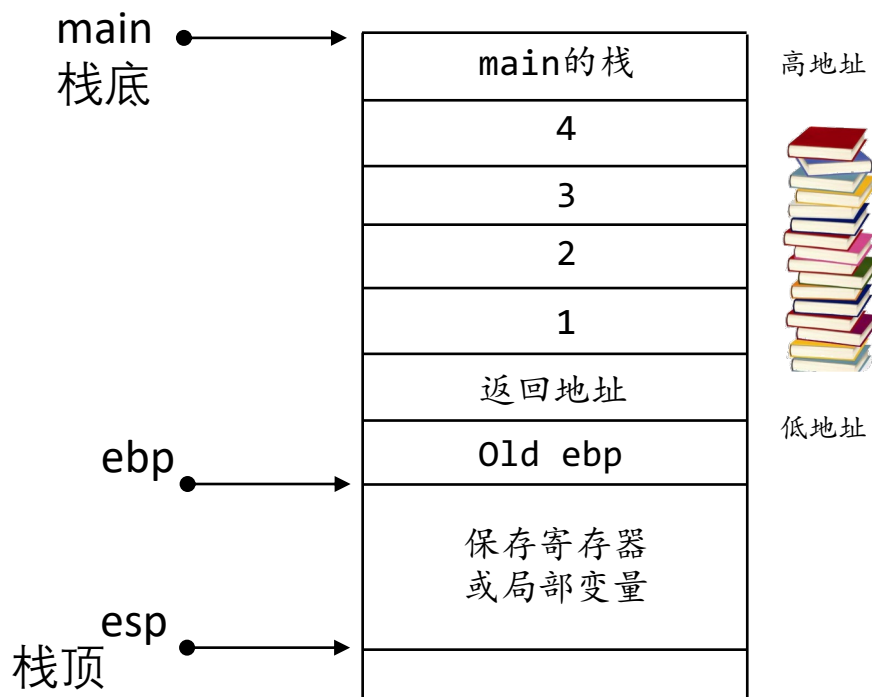
```
00000000 <foo>:  
 0: f3 0f 1e fb      endbr32  
 4: 55               push    %ebp  
 5: 89 e5            mov     %esp,%ebp  
 7: e8 fc ff ff ff   call    8 <foo+0x8>  
 c: 05 01 00 00 00   add     $0x1,%eax  
11: 8b 55 08         mov     0x8(%ebp),%edx  
14: 8b 45 0c         mov     0xc(%ebp),%eax  
17: 01 c2            add     %eax,%edx  
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1c: 01 c2            add     %eax,%edx  
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21: 01 d0            add     %edx,%eax  
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00000025 <main>:  
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31: 05 01 00 00 00   add     $0x1,%eax  
36: 6a 04            push    $0x4  
38: 6a 03            push    $0x3  
3a: 6a 02            push    $0x2  
3c: 6a 01            push    $0x1  
3e: e8 fc ff ff ff   call    3f <main+0x1a>  
43: 83 c4 10         add     $0x10,%esp  
46: b8 00 00 00 00   mov     $0x0,%eax  
4b: c9               leave   %eax  
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```

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```
int _cdecl foo(int m1, int m2, int m3, int m4);
```

```
int foo(int m1, int m2, int m3, int m4) {  
    return m1 + m2 + m3 + m4;  
}
```

```
int main() { foo(1,2,3,4);}
```



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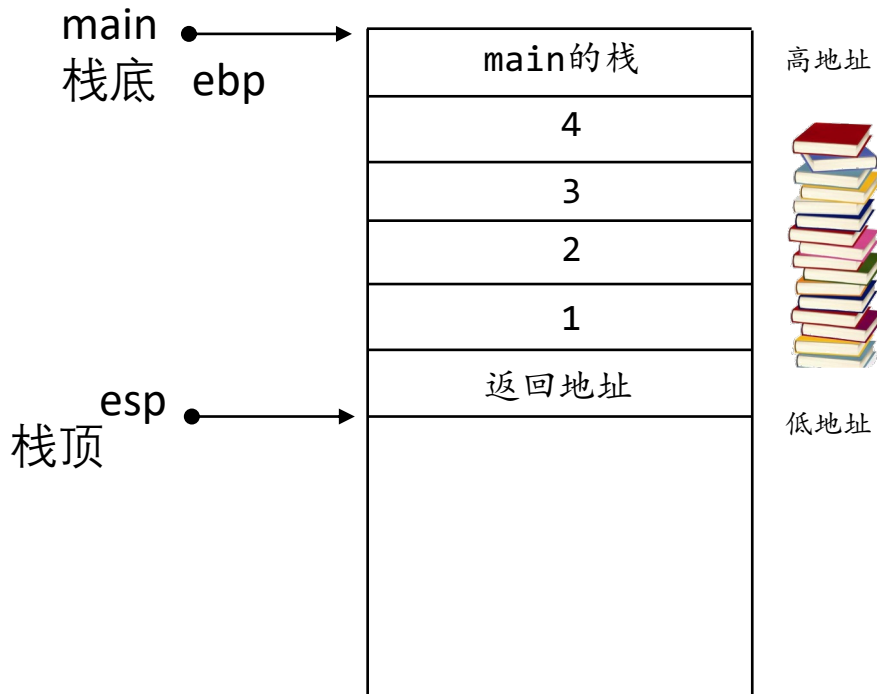
```
00000000 <foo>:  
 0: f3 0f 1e fb      endbr32  
 4: 55               push    %ebp  
 5: 89 e5            mov     %esp,%ebp  
 7: e8 fc ff ff ff   call    8 <foo+0x8>  
 c: 05 01 00 00 00   add     $0x1,%eax  
11: 8b 55 08         mov     0x8(%ebp),%edx  
14: 8b 45 0c         mov     0xc(%ebp),%eax  
17: 01 c2           add     %eax,%edx  
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1c: 01 c2           add     %eax,%edx  
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```
int foo(int m1, int m2, int m3, int m4) {  
    return m1 + m2 + m3 + m4;  
}
```

```
int main() { foo(1,2,3,4);}
```



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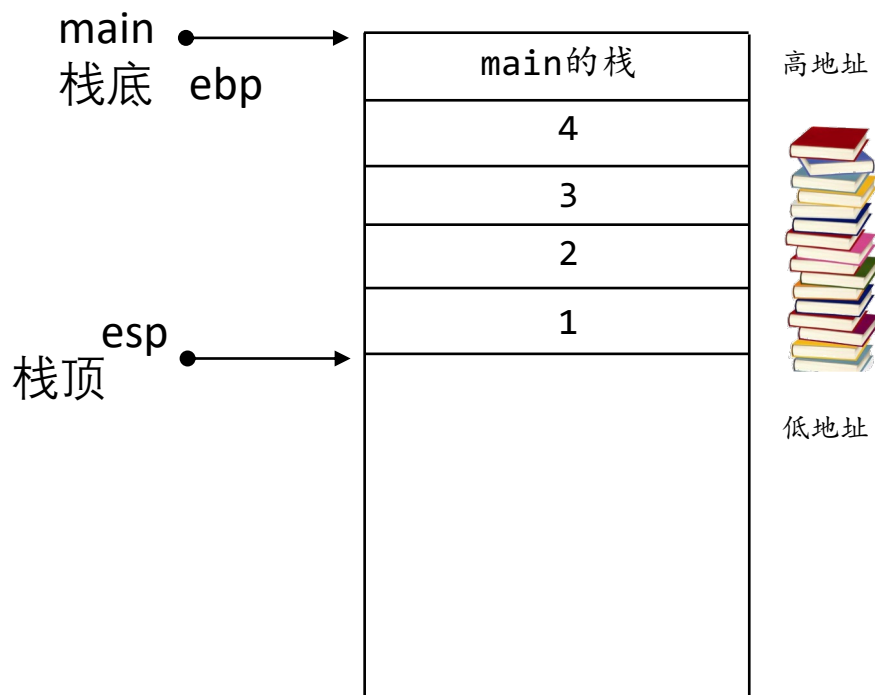
```
00000000 <foo>:  
 0: f3 0f 1e fb      endbr32  
 4: 55               push    %ebp  
 5: 89 e5            mov     %esp,%ebp  
 7: e8 fc ff ff ff   call    8 <foo+0x8>  
 c: 05 01 00 00 00   add     $0x1,%eax  
11: 8b 55 08         mov     0x8(%ebp),%edx  
14: 8b 45 0c         mov     0xc(%ebp),%eax  
17: 01 c2           add     %eax,%edx  
19: 8b 45 10         mov     0x10(%ebp),%eax  
1c: 01 c2           add     %eax,%edx  
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}
```

```
int main() { foo(1,2,3,4);}
```



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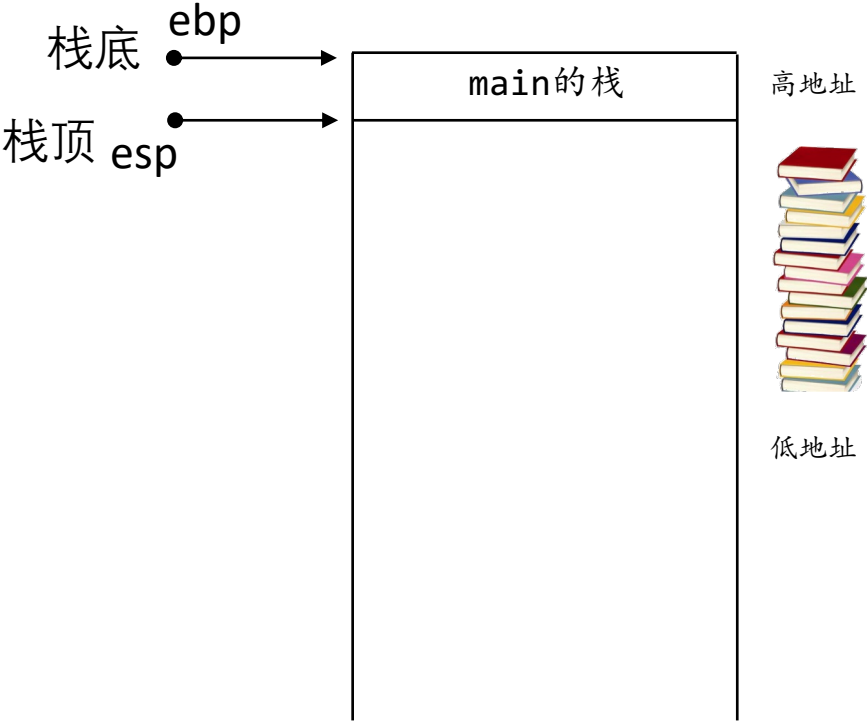
```
00000000 <foo>:  
 0: f3 0f 1e fb      endbr32  
 4: 55               push    %ebp  
 5: 89 e5            mov     %esp,%ebp  
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21: 01 d0            add     %edx,%eax  
23: 5d               pop     %ebp  
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```

```
int foo(int m1, int m2, int m3, int m4) {  
    return m1 + m2 + m3 + m4;  
}
```

```
int main() { foo(1,2,3,4);}
```



```
Disassembly of section .text:  
  
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  0:  f3 0f 1e fb      endbr32  
  4:  55               push    %ebp  
  5:  89 e5            mov     %esp,%ebp  
  7:  e8 fc ff ff ff   call   8 <foo+0x8>  
  c:  05 01 00 00 00   add     $0x1,%eax  
 11:  8b 55 08         mov     0x8(%ebp),%edx  
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 19:  8b 45 10         mov     0x10(%ebp),%eax  
 1c:  01 c2           add     %eax,%edx  
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 23:  5d              pop     %ebp  
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 29:  55               push    %ebp  
 2a:  89 e5            mov     %esp,%ebp  
 2c:  e8 fc ff ff ff   call   2d <main+0x8>  
 31:  05 01 00 00 00   add     $0x1,%eax  
 36:  6a 04            push    $0x4  
 38:  6a 03            push    $0x3  
 3a:  6a 02            push    $0x2  
 3c:  6a 01            push    $0x1  
 3e:  e8 fc ff ff ff   call   3f <main+0x1a>  
 43:  83 c4 10         add     $0x10,%esp  
 46:  b8 00 00 00 00   mov     $0x0,%eax  
 4b:  c9              leave   %eax  
 4c:  c3              ret
```

参数传递	出栈方	名字修饰
从右往左顺序压参数入栈	函数调用方	函数名前加下划线

调用惯例

```
int foo(int m1, int m2, int m3, int m4) {  
    return m1 + m2 + m3 + m4;  
}
```

	参数传递	出栈方	名字修饰
cdecl	从右往左顺序压参数入栈	函数调用方	下划线+函数名
stdcall	从右往左顺序压参数入栈	函数本身	下划线+函数名+@+参数字节数
fastcall	头两个DWORD类型或更少字节参数放入寄存器，其他参数从右往左顺序入栈	函数本身	@+函数名+@+参数字节数
pascal	从左往右顺序压参数入栈	函数本身	较复杂，见pascal文档

_foo@16

阅读汇编代码：“符号执行”

- 试着把内存/寄存器用数学符号表示出来
 - 然后假想地“单步执行”程序，用符号公式
 - James C. King. Symbolic execution and program *ACM*, 19(7), 1976.

```
0:  c3                                ret
$ gcc -c -O2 -m32 foo.c
$ objdump -d foo.o
foo.o:      file format elf32-i386

Disassembly of section .text:

00000000 <foo>:
 0:  f3 0f 1e fb                endbr32
 4:  8b 44 24 08                mov     0x8(%esp),%eax
 8:  03 44 24 04                add     0x4(%esp),%eax
 c:  03 44 24 0c                add     0xc(%esp),%eax
10:  03 44 24 10                add     0x10(%esp),%eax
14:  c3                          ret

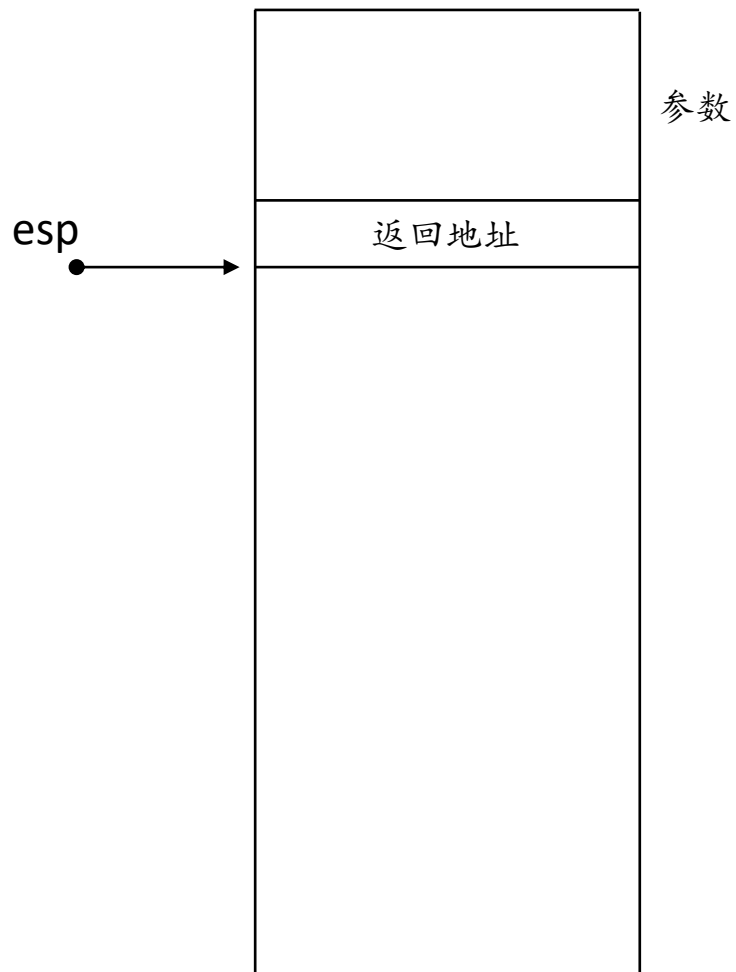
Disassembly of section .text.startup:

00000000 <main>:
 0:  f3 0f 1e fb                endbr32
 4:  31 c0                      xor     %eax,%eax
 6:  c3                          ret
```

- 编译选项：-m32 -O2 -fno-pic(便于大家理解)

```
000004f0 <foo>:
4f0:  83 ec 18                    sub     $0x18,%esp
4f3:  8d 44 24 1c                lea     0x1c(%esp),%eax
4f7:  50                          push    %eax
4f8:  e8 13 00 00 00            call    510 <bar>
4fd:  8b 44 24 20                mov     0x20(%esp),%eax
501:  83 c4 1c                    add     $0x1c,%esp
504:  c3                          ret
```

假装单步调试



000004f0 <foo>:

4f0: 83 ec 18

4f3: 8d 44 24 1c

4f7: 50

4f8: e8 13 00 00 00

4fd: 8b 44 24 20

501: 83 c4 1c

504: c3

sub \$0x18,%esp

lea 0x1c(%esp),%eax

push %eax

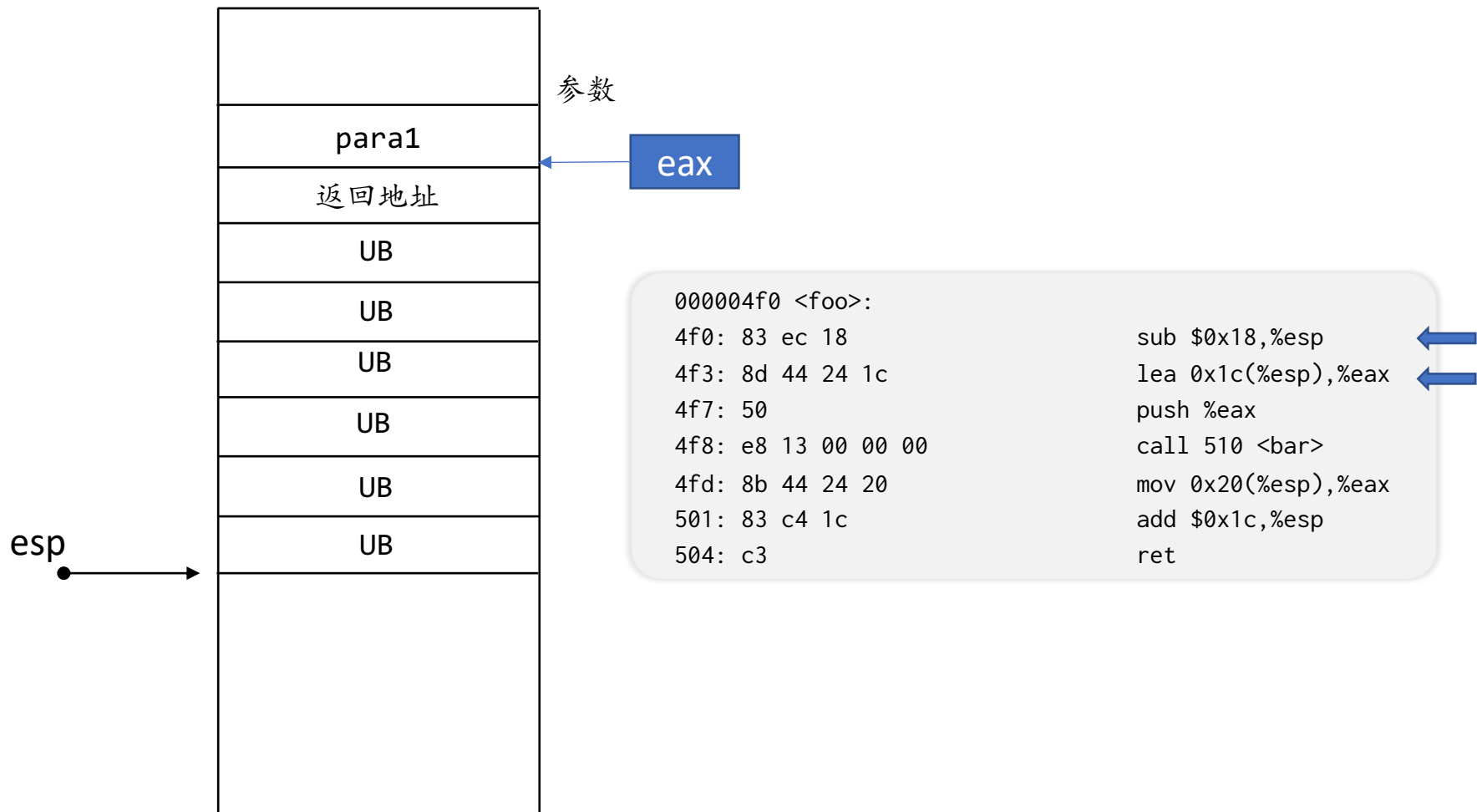
call 510 <bar>

mov 0x20(%esp),%eax

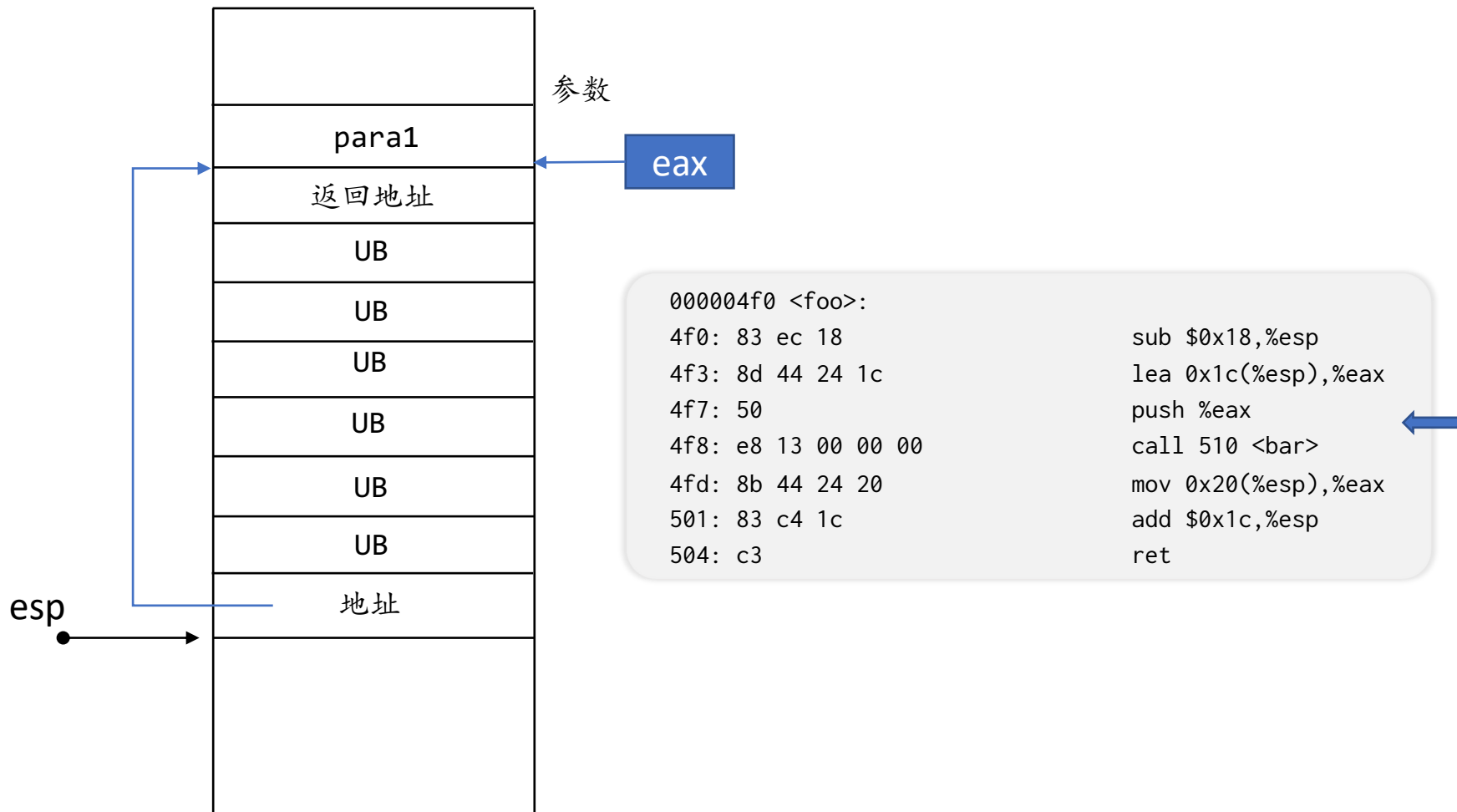
add \$0x1c,%esp

ret

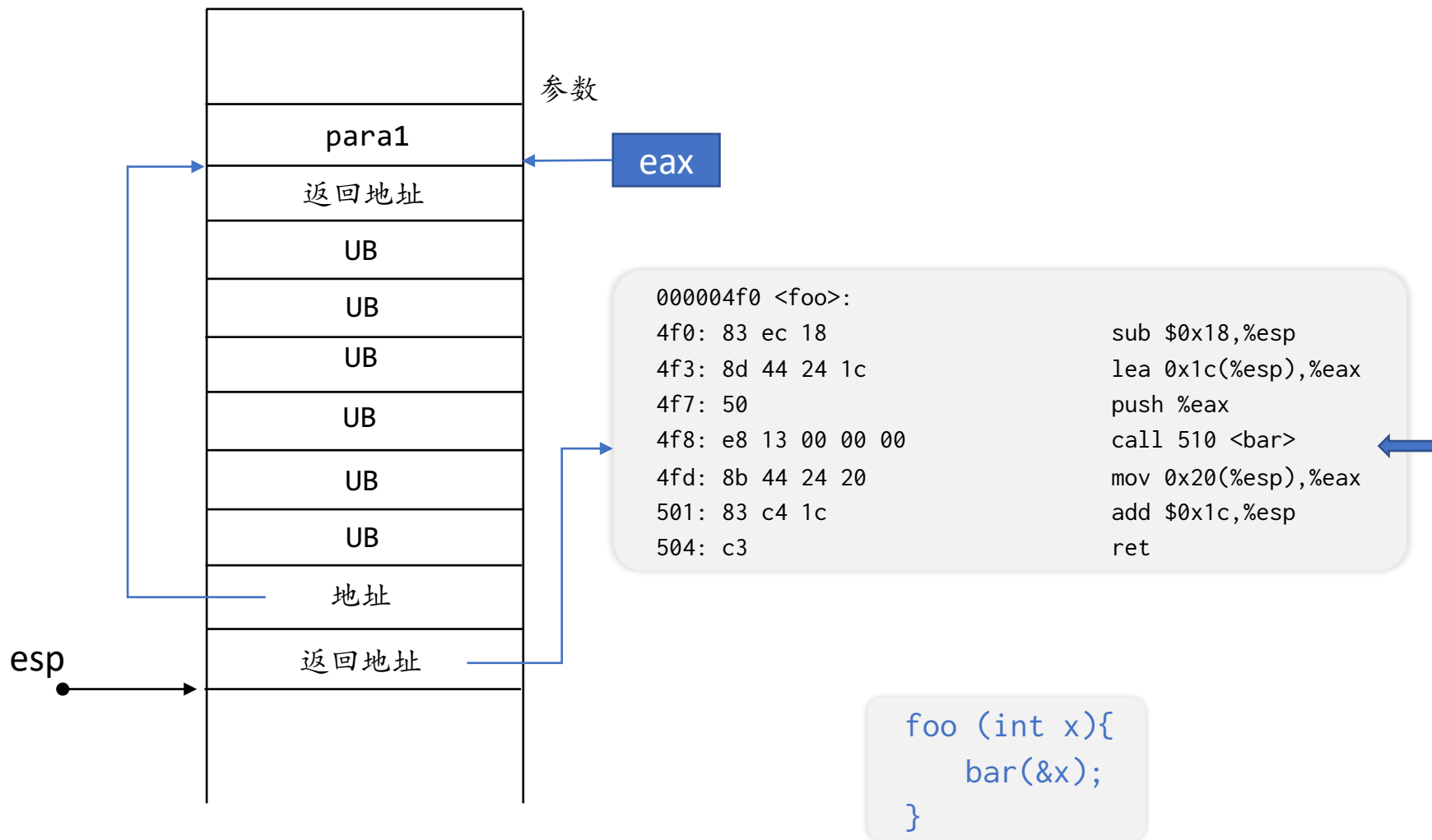
假装单步调试



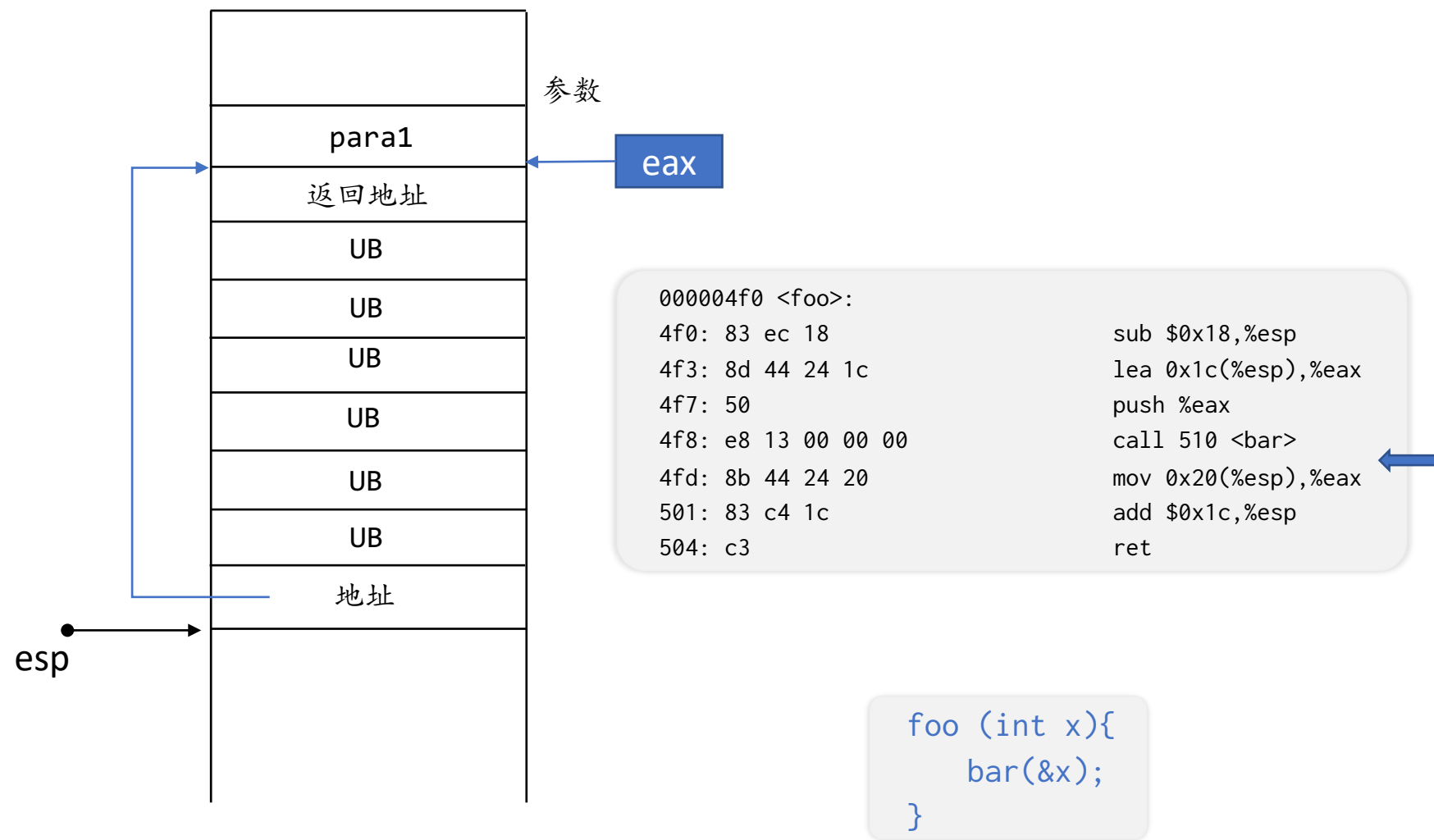
假装单步调试



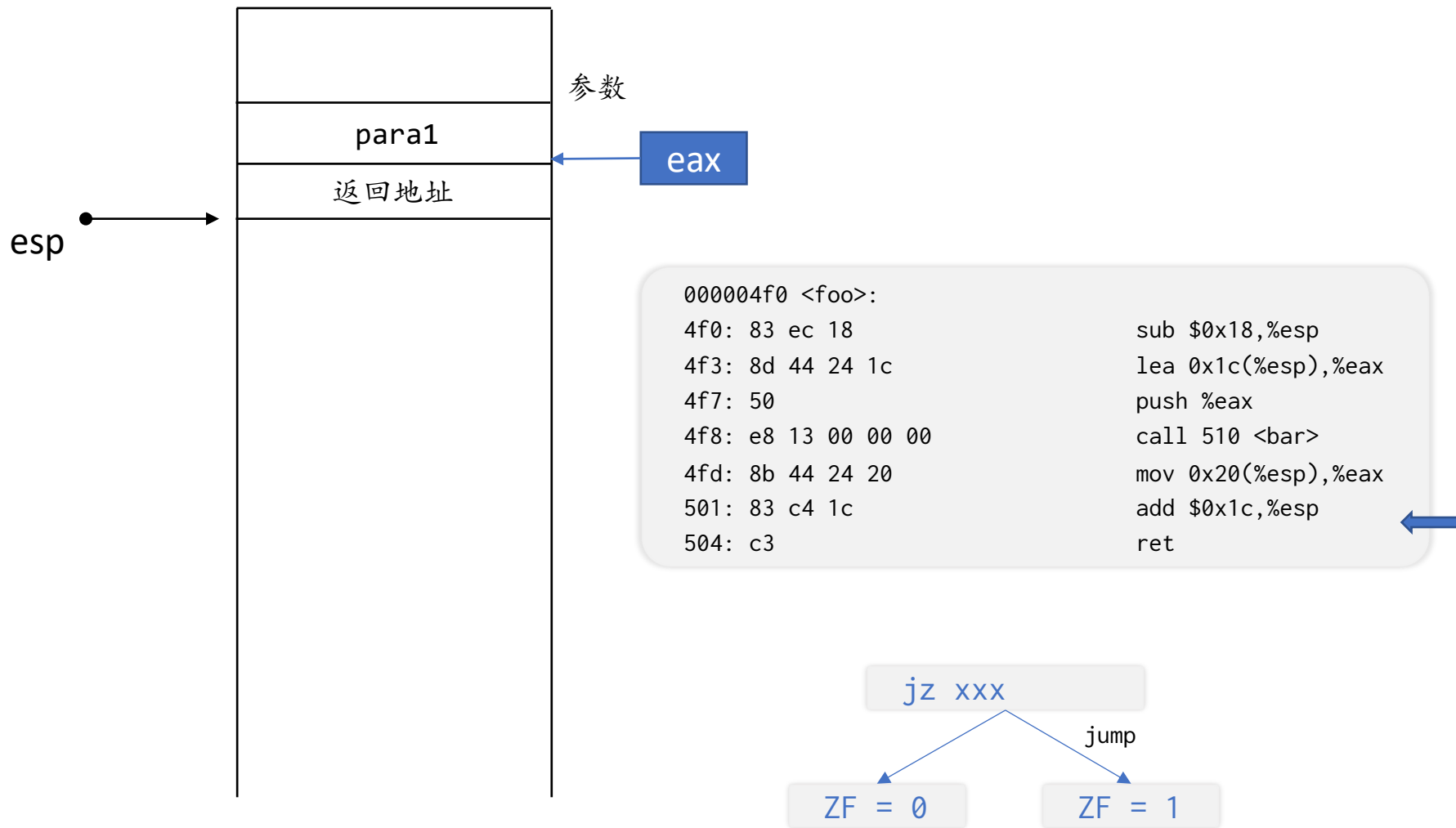
假装单步调试



假装单步调试



假装单步调试



阅读汇编代码：“符号执行”

- 试着把内存/寄存器用数学符号表示出来
 - 然后假想地“单步执行”程序，用符号公式
 - James C. King. Symbolic execution and program *ACM*, 19(7), 1976.

```
0: c3 ret
$ gcc -c -O2 -m32 foo.c
$ objdump -d foo.o
foo.o: file format elf32-i386

Disassembly of section .text:

00000000 <foo>:
 0: f3 0f 1e fb endbr32
 4: 8b 44 24 08 mov 0x8(%esp),%eax
 8: 03 44 24 04 add 0x4(%esp),%eax
 c: 03 44 24 0c add 0xc(%esp),%eax
10: 03 44 24 10 add 0x10(%esp),%eax
14: c3 ret

Disassembly of section .text.startup:

00000000 <main>:
 0: f3 0f 1e fb endbr32
 4: 31 c0 xor %eax,%eax
 6: c3 ret
```

- 编译选项：-m32 -02 -fno-pic(便于大家理解)

```
000004f0 <foo>:
4f0: 83 ec 18 sub $0x18,%esp
4f3: 8d 44 24 1c lea 0x1c(%esp),%eax
4f7: 50 push %eax
4f8: e8 13 00 00 00 call 510 <bar>
4fd: 8b 44 24 20 mov 0x20(%esp),%eax
501: 83 c4 1c add $0x1c,%esp
504: c3 ret
```

x86-64：寄存器与函数调用

寄存器 (1) : 继承自 IA32

用途	64b	低32b	低16b	低8b	8-15b
返回值	%rax	%eax	%ax	%al	%ah
调用者保存	%rbx	%ebx	%bx	%bl	%bh
参数4	%rcx	%ecx	%cx	%cl	%ch
参数3	%rdx	%edx	%dx	%dl	%dh
参数2	%rsi	%esi	%si	%sil	
参数1	%rdi	%edi	%di	%dil	
调用者保存	%rbp	%ebp	%bp	%bpl	
栈顶	%rsp	%esp	%sp	%spl	

寄存器(2): 新增加的寄存器

- x86-64 扩充了很多寄存器！
 - 于是再也不用像 IA32 那样，用堆栈传递参数了！！

用途	64b	低32b	低16b	低8b	8-15b
参数5	%r8	%r8d	%r8w	%r8b	
参数6	%r9	%r9d	%r9w	%r9b	
调用者保存	%r10	%r10d	%r10w	%r10b	
链接	%r11	%r11d	%r11w	%r11b	
C unsued	%r12	%r12d	%r12w	%r12b	
调用者保存	%r13	%r13d	%r13w	%r13b	
调用者保存	%r14	%r14d	%r14w	%r14b	
调用者保存	%r15	%r15d	%r15w	%r15b	(没有)

A + B in x86-64

```
int f(int a, int b) {  
    return a + b;  
}
```

```
00000510 <add_32>:  
  510: 8b 44 24 08      mov 0x8(%esp),%eax  
  514: 03 44 24 04      add 0x4(%esp),%eax  
  518: c3              ret
```

```
00000000000000630 <add_64>:  
  630: 8d 04 37         lea (%rdi,%rsi,1),%eax  
  633: c3              retq
```

max in x86-64

```
int max(int a, int b) {  
    if (a > b) return a;  
    else return b;  
}
```

00000514 <max_32>:

514: 8b 44 24 04	mov 0x4(%esp),%eax
518: 3b 44 24 08	cmp 0x8(%esp),%eax
51c: 7d 04	jg 522 <max+0xe>
51e: 8b 44 24 08	mov 0x8(%esp),%eax
522: c3	ret

```
movl 4(%esp), %eax  
movl 8(%esp), %edx  
cmpl %edx, %eax  
cmovl %edx, %eax  
ret
```

00000000000000640 <max_64>:

640: 39 f7	cmp %esi,%edi
642: 89 f0	mov %esi,%eax
644: 0f 4d c7	cmovge %edi,%eax
647: c3	retq

使用寄存器传递函数参数：优势

- 支持 6 个参数的传递：rdi, rsi, rdx, rcx, r8, r9
 - callee 可以随意修改这 6 个寄存器的值
 - 编译器有了更大的调度空间 (大幅减少堆栈内存访问)
- 例子：

```
void plus(int a, int b) {  
    fprintf(stdout, "%d + %d = %d\n", a, b, a + b);  
}
```

```
void plus(int a, int b) {
    fprintf(stdout, "%d + %d = %d\n", a, b, a + b);
}
```

- 实际调用的是 `__fprintf_chk@plt`
 - 需要传递的参数: `stdout, %d + %d = %d\n, a, b, a + b`
 - calling convention: `rdi, rsi, rdx, rcx, r8, r9`

0000000000000700 <plus>:

700: 44 8d 0c 37	lea (%rdi,%rsi,1),%r9d	
704: 89 f9	mov %edi,%ecx	201010 ←
706: 48 8b 3d 03 09 20 00	mov 0x200903(%rip),%rdi # XXXXXX <stdout@@GLIBC_2.2.5>	
70d: 48 8d 15 b0 00 00 00	lea 0xb0(%rip),%rdx # 7c4 <_IO_stdin_used+0x4>	
714: 41 89 f0	mov %esi,%r8d	
717: 31 c0	xor %eax,%eax	
719: be 01 00 00 00	mov \$0x1,%esi	
71e: e9 5d fe ff ff	jmpq 580 <__fprintf_chk@plt>	

rdi rsi r9 rcx

a

b

a+b

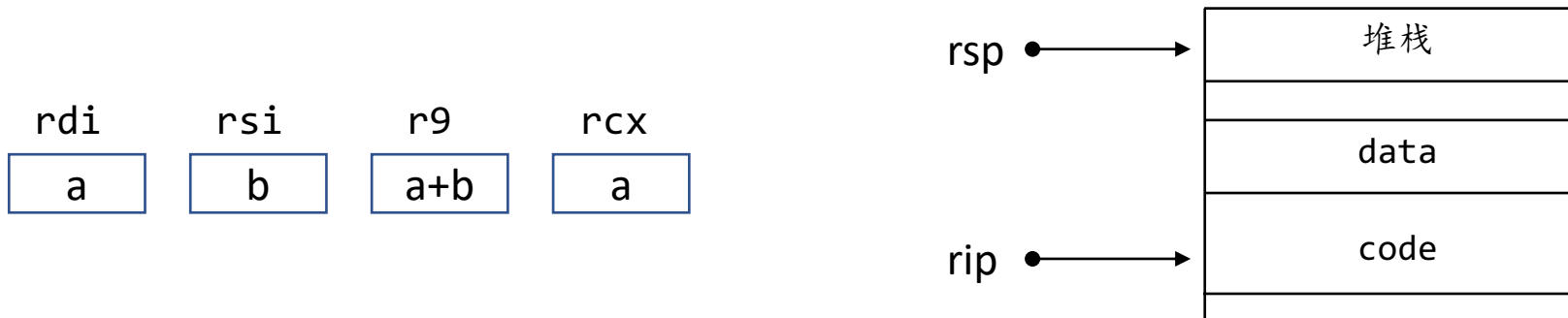
a

```
void plus(int a, int b) {
    fprintf(stdout, "%d + %d = %d\n", a, b, a + b);
}
```

- 实际调用的是 `__fprintf_chk@plt`
 - 需要传递的参数: `stdout, %d + %d = %d\n, a, b, a + b`
 - calling convention: `rdi, rsi, rdx, rcx, r8, r9`

0000000000000700 <plus>:

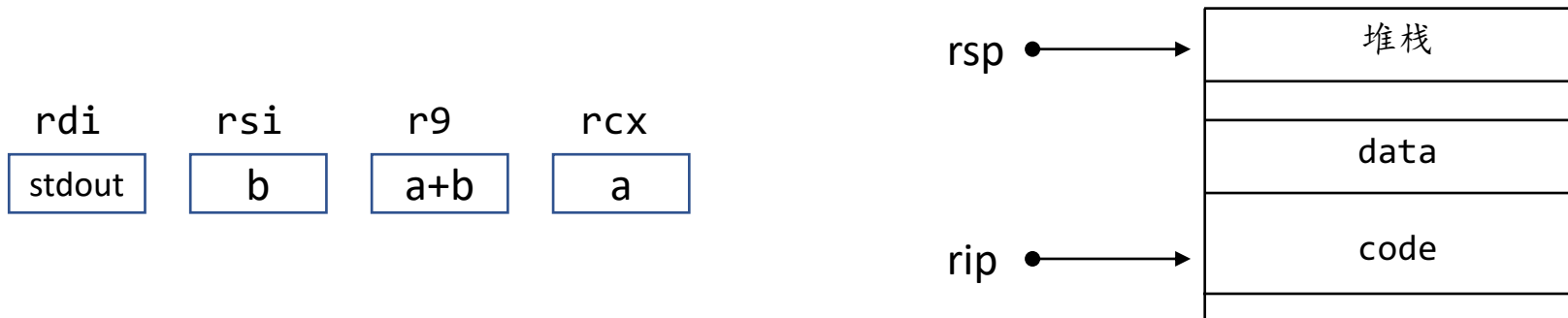
700: 44 8d 0c 37	lea (%rdi,%rsi,1),%r9d
704: 89 f9	mov %edi,%ecx
706: 48 8b 3d 03 09 20 00	mov 0x200903(%rip),%rdi # 201010 <stdout@@GLIBC_2.2.5>
70d: 48 8d 15 b0 00 00 00	lea 0xb0(%rip),%rdx # 7c4 <_IO_stdin_used+0x4>
714: 41 89 f0	mov %esi,%r8d
717: 31 c0	xor %eax,%eax
719: be 01 00 00 00	mov \$0x1,%esi
71e: e9 5d fe ff ff	jmpq 580 <__fprintf_chk@plt>



```
void plus(int a, int b) {
    fprintf(stdout, "%d + %d = %d\n", a, b, a + b);
}
```

- 实际调用的是 `__fprintf_chk@plt`
 - 需要传递的参数: `stdout, %d + %d = %d\n, a, b, a + b`
 - calling convention: `rdi, rsi, rdx, rcx, r8, r9`

```
0000000000000700 <plus>:
700: 44 8d 0c 37          lea (%rdi,%rsi,1),%r9d
704: 89 f9               mov %edi,%ecx
706: 48 8b 3d 03 09 20 00  mov 0x200903(%rip),%rdi # 201010 <stdout@@GLIBC_2.2.5>
70d: 48 8d 15 b0 00 00 00  lea 0xb0(%rip),%rdx # 7c4 <_IO_stdin_used+0x4>
714: 41 89 f0            mov %esi,%r8d
717: 31 c0              xor %eax,%eax
719: be 01 00 00 00      mov $0x1,%esi
71e: e9 5d fe ff ff      jmpq 580 <__fprintf_chk@plt>
```

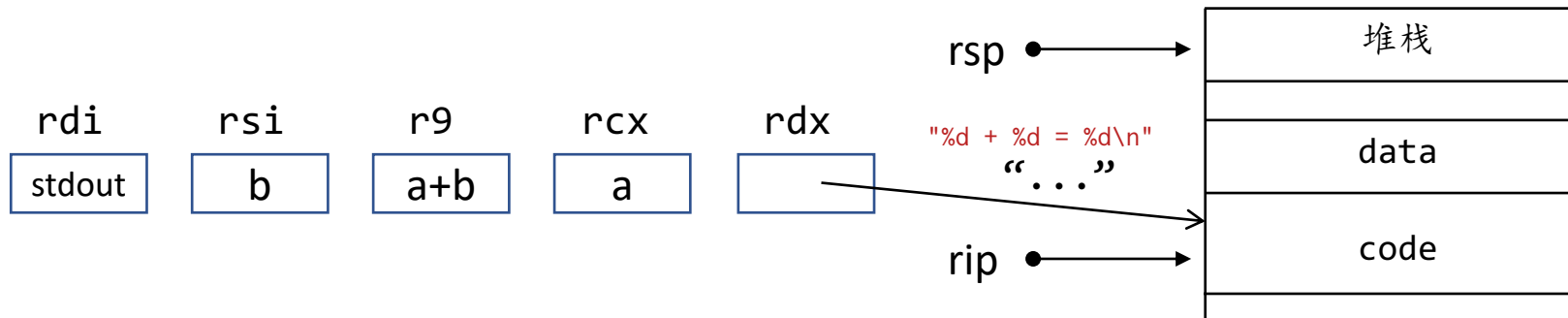


```
void plus(int a, int b) {
    fprintf(stdout, "%d + %d = %d\n", a, b, a + b);
}
```

- 实际调用的是 `__fprintf_chk@plt`
 - 需要传递的参数：`stdout`, `%d + %d = %d\n`, `a`, `b`, `a + b`
 - calling convention: `rdi`, `rsi`, `rdx`, `rcx`, `r8`, `r9`

0000000000000700 <plus>:

700: 44 8d 0c 37	lea (%rdi,%rsi,1),%r9d
704: 89 f9	mov %edi,%ecx
706: 48 8b 3d 03 09 20 00	mov 0x200903(%rip),%rdi # 201010 <stdout@@GLIBC_2.2.5>
70d: 48 8d 15 b0 00 00 00	lea 0xb0(%rip),%rdx # 7c4 <_IO_stdin_used+0x4>
714: 41 89 f0	mov %esi,%r8d
717: 31 c0	xor %eax,%eax
719: be 01 00 00 00	mov \$0x1,%esi
71e: e9 5d fe ff ff	jmpq 580 <__fprintf_chk@plt>

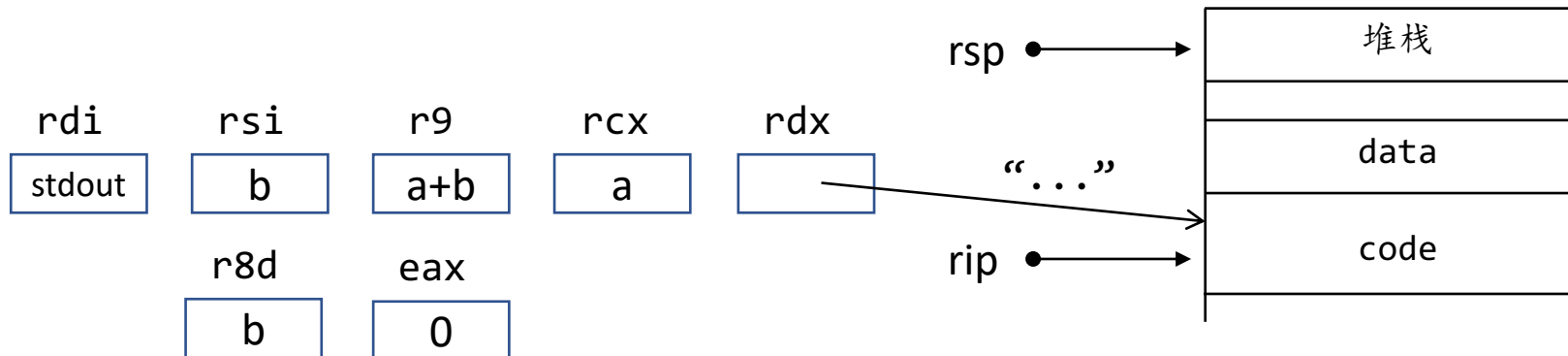


```
void plus(int a, int b) {
    fprintf(stdout, "%d + %d = %d\n", a, b, a + b);
}
```

- 实际调用的是 `__fprintf_chk@plt`
 - 需要传递的参数：`stdout`, `%d + %d = %d\n`, `a`, `b`, `a + b`
 - calling convention: `rdi`, `rsi`, `rdx`, `rcx`, `r8`, `r9`

0000000000000700 <plus>:

700: 44 8d 0c 37	lea (%rdi,%rsi,1),%r9d
704: 89 f9	mov %edi,%ecx
706: 48 8b 3d 03 09 20 00	mov 0x200903(%rip),%rdi # 201010 <stdout@@GLIBC_2.2.5>
70d: 48 8d 15 b0 00 00 00	lea 0xb0(%rip),%rdx # 7c4 <_IO_stdin_used+0x4>
714: 41 89 f0	mov %esi,%r8d
717: 31 c0	xor %eax,%eax
719: be 01 00 00 00	mov \$0x1,%esi
71e: e9 5d fe ff ff	jmpq 580 <__fprintf_chk@plt>

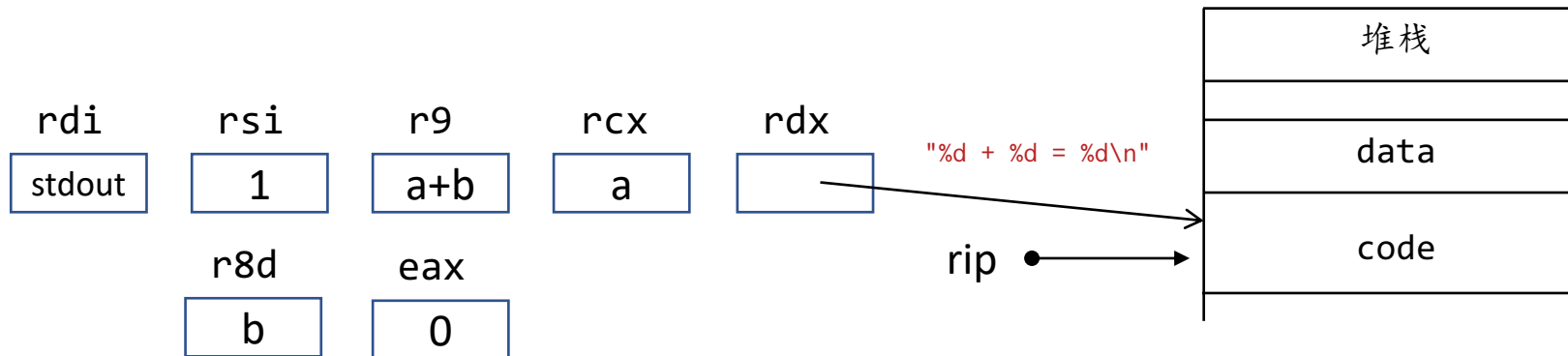


```
void plus(int a, int b) {
    fprintf(stdout, "%d + %d = %d\n", a, b, a + b);
}
```

- 实际调用的是 `__fprintf_chk@plt`
 - 需要传递的参数: `stdout, %d + %d = %d\n, a, b, a + b`
 - calling convention: `rdi, rsi, rdx, rcx, r8, r9`

0000000000000700 <plus>:

700: 44 8d 0c 37	lea (%rdi,%rsi,1),%r9d
704: 89 f9	mov %edi,%ecx
706: 48 8b 3d 03 09 20 00	mov 0x200903(%rip),%rdi # 201010 <stdout@@GLIBC_2.2.5>
70d: 48 8d 15 b0 00 00 00	lea 0xb0(%rip),%rdx # 7c4 <_IO_stdin_used+0x4>
714: 41 89 f0	mov %esi,%r8d
717: 31 c0	xor %eax,%eax
719: be 01 00 00 00	mov \$0x1,%esi ←
71e: e9 5d fe ff ff	jmpq 580 <__fprintf_chk@plt>



一些更多的分析

- plus的最后一条指令

```
71e: e9 5d fe ff ff      jmpq 580 <__fprintf_chk@plt>
```

- 并不是调用的printf，而是调用的[有堆栈检查的版本](#)
 - 准备参数时有 `mov $0x1,%esi`
- 直接jmp是因为函数末尾有一条ret指令
 - 借用了__fprintf_chk@plt的ret指令返回到plus的调用者
 - 如果有返回值，就会生成call指令；如果plus返回printf的结果，依然是jmp
 - 省的不止是一条指令
 - 连续的ret对分支预测是很大的挑战

对比32位printf

- 好读，但指令执行起来会稍慢一些

```
void plus(int a, int b) {  
    fprintf(stdout, "%d + %d = %d\n", a, b, a + b);  
}
```

000005b4 <plus>:

5b4: 83 ec 14	sub	\$0x14,%esp
5b7: 8b 44 24 18	mov	0x18(%esp),%eax
5bb: 8b 54 24 1c	mov	0x1c(%esp),%edx
5bf: 8d 0c 10	lea	(%eax,%edx,1),%ecx
5c2: 51	push	%ecx
5c3: 52	push	%edx
5c4: 50	push	%eax
5c5: 68 60 06 00 00	push	\$0x660
5ca: 6a 01	push	\$0x1
5cc: ff 35 00 00 00 00	pushl	0x0
5d2: e8 fc ff ff ff	call	5d3 <plus+0x1f>
5d7: 83 c4 2c	add	\$0x2c,%esp
5da: c3	ret	
5db: 90	nop	

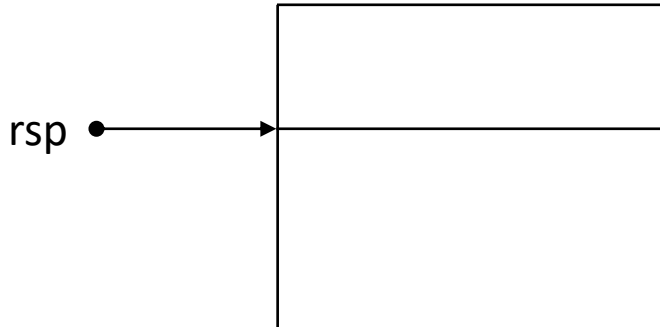
一个有趣的小问题

- x86-64 按寄存器传递参数
 - `void f(int x) {... &x ...}` 会发生什么？
- 编译器会给参数分配内存，保证后续访问合法
 - 给编译器带来了轻微的负担
 - 但编译器并不觉得这是负担

```
void bar(int *);  
int foo(int x) {  
    bar(&x);  
    return x;  
}
```

假装单步调试

```
void bar(int *);  
int foo(int x) {  
    bar(&x);  
    return x;  
}
```



```
$ gcc -c -O2 foo.c  
$ objdump -d foo.o
```

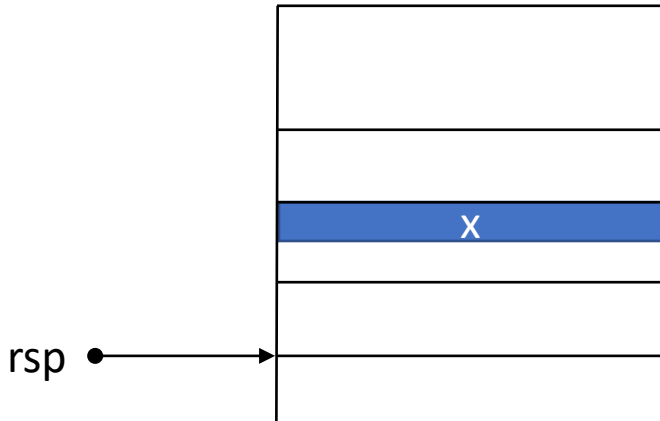
```
foo.o:      file format elf64-x86-64
```

```
Disassembly of section .text:
```

```
0000000000000000 <foo>:  
  0:  f3 0f 1e fa                endbr64  
  4:  48 83 ec 18                sub    $0x18,%rsp  
  8:  89 7c 24 0c                mov    %edi,0xc(%rsp)  
  c:  48 8d 7c 24 0c            lea    0xc(%rsp),%rdi  
 11:  e8 00 00 00 00            call   16 <foo+0x16>  
 16:  8b 44 24 0c                mov    0xc(%rsp),%eax  
 1a:  48 83 c4 18                add    $0x18,%rsp  
 1e:  c3                        ret
```

假装单步调试

```
void bar(int *);  
int foo(int x) {  
    bar(&x);  
    return x;  
}
```



```
$ gcc -c -O2 foo.c  
$ objdump -d foo.o  
  
foo.o:      file format elf64-x86-64
```

Disassembly of section `.text`:

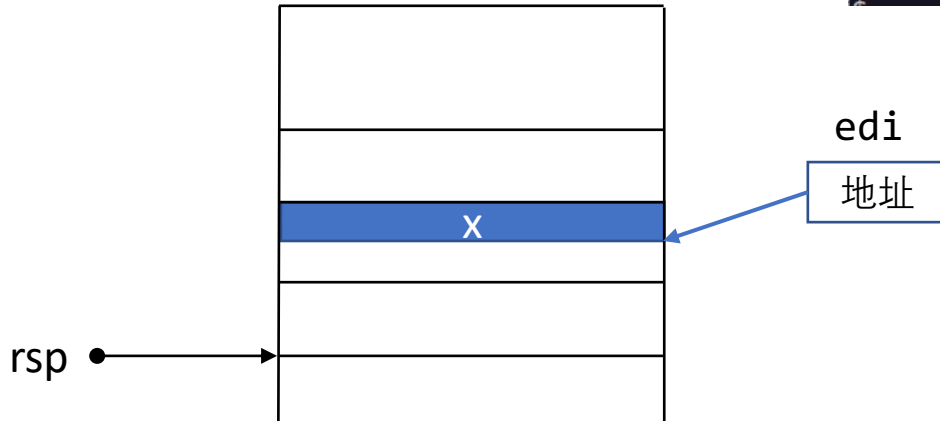
```
0000000000000000 <foo>:  
0:  f3 0f 1e fa      endbr64  
4:  48 83 ec 18      sub    $0x18,%rsp  
8:  89 7c 24 0c      mov    %edi,0xc(%rsp)  
c:  48 8d 7c 24 0c   lea    0xc(%rsp),%rdi  
11: e8 00 00 00 00   call   16 <foo+0x16>  
16: 8b 44 24 0c      mov    0xc(%rsp),%eax  
1a: 48 83 c4 18      add    $0x18,%rsp  
1e:  c3              ret
```

edi

x

假装单步调试

```
void bar(int *);  
int foo(int x) {  
    bar(&x);  
    return x;  
}
```



```
$ gcc -c -O2 foo.c  
$ objdump -d foo.o  
  
foo.o:      file format elf64-x86-64
```

Disassembly of section .text:

```
0000000000000000 <foo>:  
  0:  f3 0f 1e fa                endbr64  
  4:  48 83 ec 18                sub     $0x18,%rsp  
  8:  89 7c 24 0c                mov     %edi,0xc(%rsp)  
  c:  48 8d 7c 24 0c            lea     0xc(%rsp),%rdi  
 11:  e8 00 00 00 00            call    16 <foo+0x16>  
 16:  8b 44 24 0c                mov     0xc(%rsp),%eax  
 1a:  48 83 c4 18                add     $0x18,%rsp  
 1e:  c3                        ret
```

x86-64程序：更多的例子

swap in x86-64

- 总的来说，x86-64是更现代的体系结构，更精简的指令序列
 - `void swap(int *x, int *y);` 交换两个指针指向的数字

```
mov    0x8(%esp),%edx
mov    0xc(%esp),%eax
mov    (%edx),%ecx
mov    (%eax),%ebx
mov    %ebx, (%edx)
mov    %ecx, (%eax)
pop    %ebx
```

```
mov    (%rdi),%eax
mov    (%rsi),%edx
mov    %edx, (%rdi)
mov    %eax, (%rsi)
```

例子：循环

```
int fact(int n) {  
    int res = 1;  
    do { res *= n; n--; } while (n > 0);  
    return res;  
}
```

```
        mov     $0x1,%eax  
        nopl    (%rax)  
.L1:    imul    %edi,%eax  
        sub     $0x1,%edi  
        test    %edi,%edi  
        jg      .L1  
        repz    retq
```

- 两个诡异代码：
 - `nopl (%rax)`：内存对齐 (padding)
 - `repz retq`：防止连续分支指令

例子：递归

0000000000000704 <fib>:

```
704: 55          push    %rbp
705: 53          push    %rbx
706: 89 fd       mov     %edi,%ebp
708: 31 db       xor     %ebx,%ebx
70a: 48 83 ec 08  sub     $0x8,%rsp
70e: 83 fd 01     cmp     $0x1,%ebp
711: 7e 0f       jle     722 <fib+0x1e>
713: 8d 7d ff     lea     -0x1(%rbp),%edi
716: 83 ed 02     sub     $0x2,%ebp
719: e8 e6 ff ff  callq   704 <fib>
71e: 01 c3       add     %eax,%ebx
720: eb ec       jmp     70e <fib+0xa>
722: 8d 43 01     lea     0x1(%rbx),%eax
725: 5a         pop     %rdx
726: 5b         pop     %rbx
727: 5d         pop     %rbp
728: c3         retq
```

Inline Assembly

在 C 代码中嵌入汇编

- 编译器：把 C 代码 “原样翻译” 成汇编代码
 - 那我们是不是可以在语句之间插入汇编呢？
 - 可以！编译器就做个 “复制粘贴”

```
int foo(int x, int y) {  
    x++; y++;  
    asm ("endbr64;"  
        ".byte 0xf3, 0x0f, 0x1e, 0xfa"  
    );  
    return x + y;  
}
```

```
$ objdump -d foo.o  
  
foo.o:      file format elf64-x86-64  
  
Disassembly of section .text:  
  
0000000000000000 <foo>:  
0:  f3 0f 1e fa      endbr64  
4:  55               push    %rbp  
5:  48 89 e5         mov     %rsp,%rbp  
8:  89 7d fc         mov     %edi,-0x4(%rbp)  
b:  89 75 f8         mov     %esi,-0x8(%rbp)  
e:  83 45 fc 01      addl    $0x1,-0x4(%rbp)  
12: 83 45 f8 01      addl    $0x1,-0x8(%rbp)  
16: f3 0f 1e fa      endbr64  
1a: f3 0f 1e fa      endbr64  
1e: 8b 55 fc         mov     -0x4(%rbp),%edx  
21: 8b 45 f8         mov     -0x8(%rbp),%eax  
24: 01 d0           add     %edx,%eax  
26: 5d              pop     %rbp  
27: c3              ret
```

在汇编中访问 C 世界的表达式

```
int foo(int x, int y) {  
    int z;  
    x++; y++;  
    asm (  
        "addl %1, %2; "  
        "movl %2, %0; "  
        : "=r"(z) // output  
        : "r"(x), "r"(y) // input  
    );  
    return z;  
}
```

gcc -O0

```
0000000000000000 <foo>:  
  0:  f3 0f 1e fa      endbr64  
  4:  55               push    %rbp  
  5:  48 89 e5         mov     %rsp,%rbp  
  8:  89 7d ec         mov     %edi,-0x14(%rbp)  
  b:  89 75 e8         mov     %esi,-0x18(%rbp)  
  e:  83 45 ec 01      addl    $0x1,-0x14(%rbp)  
 12:  83 45 e8 01      addl    $0x1,-0x18(%rbp)  
 16:  8b 45 ec         mov     -0x14(%rbp),%eax  
 19:  8b 55 e8         mov     -0x18(%rbp),%edx  
 1c:  01 c2           add     %eax,%edx  
 1e:  89 d0           mov     %edx,%eax  
 20:  89 45 fc         mov     %eax,-0x4(%rbp)  
 23:  8b 45 fc         mov     -0x4(%rbp),%eax  
 26:  5d             pop     %rbp  
 27:  c3             ret
```

在汇编中访问 C 世界的表达式

```
int foo(int x, int y) {  
    int z;  
    x++; y++;  
    asm (  
        "addl %1, %2; "  
        "movl %2, %0; "  
        : "=r"(z) // output  
        : "r"(x), "r"(y) // input  
    );  
    return z;  
}
```

gcc -O2

```
0000000000000000 <foo>:  
0:  f3 0f 1e fa      endbr64  
4:  83 c7 01          add    $0x1,%edi  
7:  8d 46 01          lea    0x1(%rsi),%eax  
a:  01 f8            add    %edi,%eax  
c:  89 c0            mov    %eax,%eax  
e:  c3              ret
```

在汇编中访问 C 世界的表达式

```
int foo(int x, int y) {  
    int z;  
    x++; y++;  
    asm (  
        "addl %1, %2; "  
        "movl %2, %0; "  
        : "=&r"(z) // output  
        : "r"(x), "r"(y) // input  
    );  
    return z;  
}
```

gcc -O2

```
0000000000000000 <foo>:  
0:  f3 0f 1e fa          endbr64  
4:  83 c7 01             add    $0x1,%edi  
7:  83 c6 01             add    $0x1,%esi  
a:  01 fe               add    %edi,%esi  
c:  89 f0               mov    %esi,%eax  
e:  c3                  ret
```

与编译器交互

- 实际的编译器可能会
 - 将某个变量分配给某个寄存器
 - inline assembly 改写寄存器就会导致错误
 - → clobber list

```
asm asm-qualifiers ( AssemblerTemplate
                     : OutputOperands
                     [ : InputOperands
                     [ : Clobbers ] ] )
```

```
asm asm-qualifiers ( AssemblerTemplate
                     : OutputOperands
                     : InputOperands
                     : Clobbers
                     : GotoLabels )
```

```

1 #include <stdio.h>
2
3 int inc1(int src) {
4     int dst;
5
6     asm("mov %1, %0\n\t"
7         "add $1, %0"
8         : "=r"(dst)
9         : "r"(src));
10
11     return dst;
12 }
13
14 int inc2(int src) {
15     int dst;
16
17     asm("mov %1, %0\n\t"
18         "mov $3, %%eax\n\t"
19         "add $1, %0"
20         : "=r"(dst)
21         : "r"(src));
22
23     return dst;
24 }

```

```

26 int inc3(int src) {
27     int dst;
28
29     asm("mov %1, %0\n\t"
30         "mov $3, %%eax\n\t"
31         "add $1, %0"
32         : "=r"(dst)
33         : "r"(src)
34         : "%eax");
35
36     return dst;
37 }
38
39 int main(int argc, char **argv) {
40     printf("inc1: %d\n", inc1(2));
41     printf("inc2: %d\n", inc2(4));
42     printf("inc3: %d\n", inc3(2));
43     return 0;
44 }

```

```

inc1: 2
inc2: 4
inc3: 2

```

```

endbr64
push    %rbp
mov     %rsp,%rbp
mov     %edi,-0x14(%rbp)
mov     -0x14(%rbp),%eax
mov     %eax,%eax
mov     $0x3,%eax
add     $0x1,%eax
mov     %eax,-0x4(%rbp)
mov     -0x4(%rbp),%eax
pop     %rbp
ret

```

与编译器交互

- 实际的编译器可能会
 - 代码优化
 - 例如 assembly 的返回值没有用到，就可以删除
 - → **volatile**

```
int foo(int x, int y) {  
    int z;  
    x++; y++;  
    asm (  
        "addl %1, %2; "  
        "movl %2, %0; "  
        : "=r"(z) // output  
        : "r"(x), "r"(y) // input  
    );  
    return 0;  
}
```

gcc -O2

```
0000000000000000 <foo>:  
0:  f3 0f 1e fa      endbr64  
4:  31 c0            xor    %eax,%eax  
6:  c3              ret
```

在汇编中访问 C 世界的表达式

```
int foo(int x, int y) {  
    int z;  
    x++; y++;  
    asm volatile(  
        "addl %1, %2; "  
        "movl %2, %0; "  
        : "=r"(z) // output  
        : "r"(x), "r"(y) // input  
    );  
    return 0;  
}
```

gcc -O2

```
0000000000000000 <foo>:  
0: f3 0f 1e fa          endbr64  
4: 83 c7 01             add    $0x1,%edi  
7: 83 c6 01             add    $0x1,%esi  
a: 01 fe             add    %edi,%esi  
c: 89 f7             mov    %esi,%edi  
e: 31 c0             xor    %eax,%eax  
10: c3             ret
```

总结

对汇编感到很痛苦？

- 两个建议
 - 原理：想一想 YEMU 和 NEMU
 - 实践：“模拟调试”
 - 指令不过是 CPU 执行的基本单元
- 同时，付出也是必要的
 - 但你一旦掌握了分析汇编代码的方法
 - x86-64 也不可怕
 - AArch64 也不可怕
 - 学会查手册

End. (你找对手册了吗？)

[x86-64 Machine-Level Programming](#)
[How to Use Inline Assembly Language in C Code](#)